

# Quiz

Let  $f, g : A \rightarrow B$  be two arrows. Define the *equalizer* of  $f$  and  $g$ .

# Definition

A map  $e$  which is an equalizer of two arrows is sometimes called a *regular* arrow.

# Lemma

Every regular arrow is strong.

# Proof

Let  $e$  be a regular arrow and

$$E \xrightarrow{e} A \begin{array}{c} \xrightarrow{f} \\ \xrightarrow{g} \end{array} B$$

be its equalizer diagram.

Let  $t : T \rightarrow A$  be an arbitrary arrow and


$$\begin{array}{ccc} P & \xrightarrow{\tilde{t}} & E \\ \tilde{e} \downarrow & & \downarrow e \\ T & \xrightarrow{t} & A \end{array}$$

be a pullback diagram. Suppose that  $\tilde{e}$  is epic.

Consider now the diagram

$$\begin{array}{ccccc} P & \xrightarrow{\tilde{t}} & A & & \\ \tilde{e} \downarrow & & \downarrow e & & \\ T & \xrightarrow{t} & A & \xrightarrow{f} & B \\ & & & \xrightarrow{g} & \end{array}$$

*It commutes* (with the exception that  $f \neq g$ ). This means


$$f \circ t \circ \tilde{e} = g \circ t \circ \tilde{e}$$

- $f \circ t \circ \tilde{e} = g \circ t \circ \tilde{e}$

- $f \circ t = g \circ t$  (since  $\tilde{e}$  is epic) which implies that

- $f \circ t \circ \tilde{e} = g \circ t \circ \tilde{e}$

- $f \circ t = g \circ t$  (since  $\tilde{e}$  is epic) which implies that

- the diagram

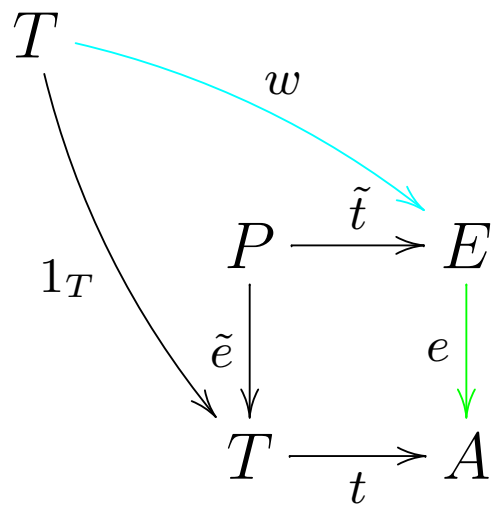
$$\begin{array}{ccccc} E & \xrightarrow{e} & A & \begin{array}{l} \xrightarrow{f} \\ \xrightarrow{g} \end{array} & B \\ & & \uparrow t & & \\ & & T & & \end{array}$$

commutes.

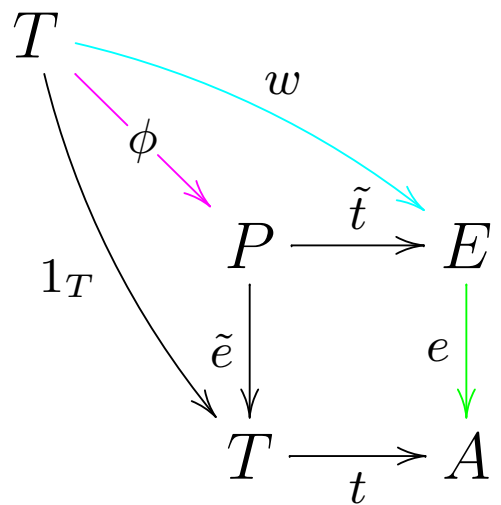
Thus, there exists a unique  $w$  so that

$$\begin{array}{ccccc} E & \xrightarrow{e} & A & \begin{array}{l} \xrightarrow{f} \\ \xrightarrow{g} \end{array} & B \\ & \swarrow w & \uparrow t & & \\ & & T & & \end{array}$$

commutes. Thus, so too does



Thus, there exists a unique  $\phi$  so that



commutes.

Now,

•  $\tilde{e} \circ \phi = 1_T$

Now,

- $\tilde{e} \circ \phi = 1_T$

- $\tilde{e} \circ \phi \circ \tilde{e} = \tilde{e} = \tilde{e}1_P$

Now,

- $\tilde{e} \circ \phi = 1_T$

- $\tilde{e} \circ \phi \circ \tilde{e} = \tilde{e} = \tilde{e}1_P$

- $\phi \circ \tilde{e} = 1_P$  (since  $\tilde{e}$ , being the pullback of the monic  $e$  is itself monic)

Now,

- $\tilde{e} \circ \phi = 1_T$
- $\tilde{e} \circ \phi \circ \tilde{e} = \tilde{e} = \tilde{e}1_P$
- $\phi \circ \tilde{e} = 1_P$  (since  $\tilde{e}$ , being the pullback of the monic  $e$  is itself monic)
- Thus, we have shown that if the pullback of  $e$  is epi it must also be an isomorphism. This is the definition of a strong arrow.

# Lemma

Let  $X \begin{array}{c} \xrightarrow{f} \\ \xrightarrow{g} \end{array} Y$  be two functions. Then the equalizer of  $f$  and  $g$  is the set  $E := \{x \in X \mid f(x) = g(x)\}$  and the function  $e : E \rightarrow X$  defined by  $e(w) = w$  - that is, the inclusion.